

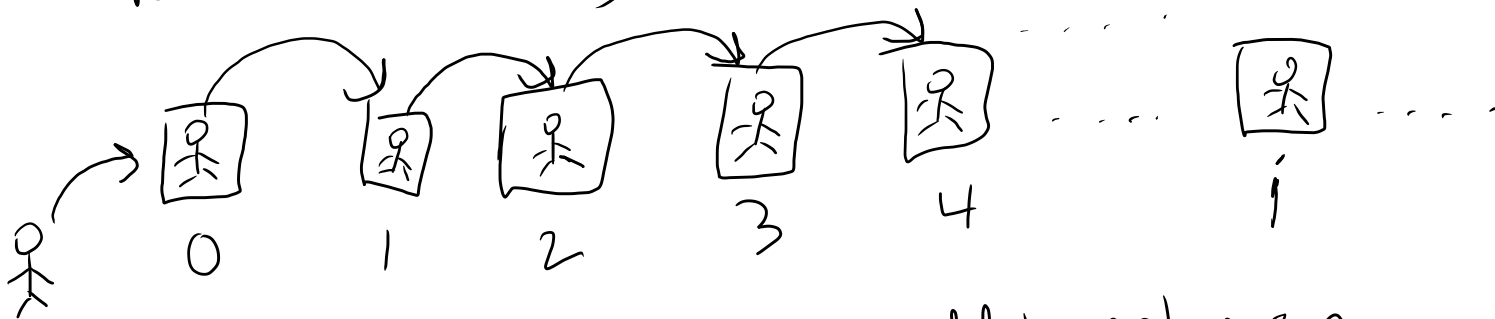
Infinity

- Lets say you have a hotel with finitely many rooms, e.g. 5 rooms.



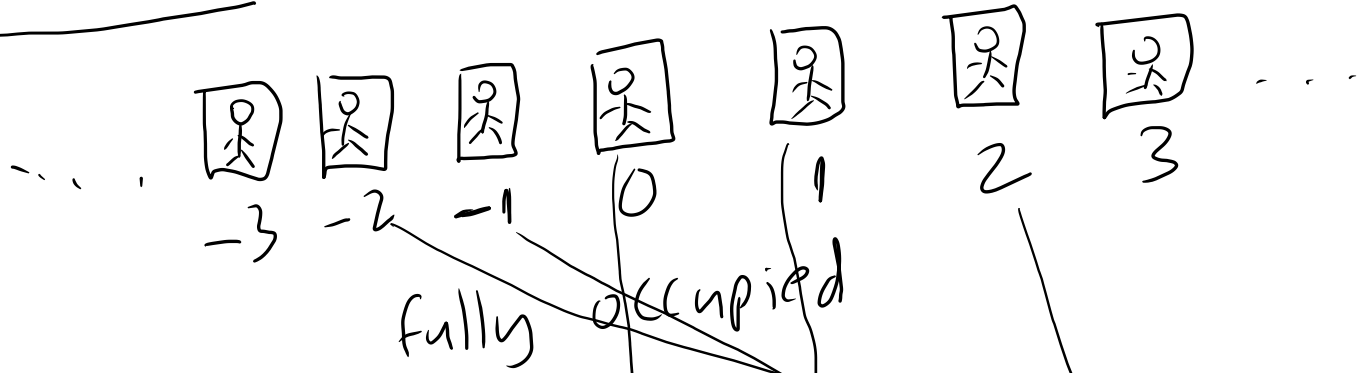
- Suppose all rooms are occupied
- Suppose one more person shows up, demands a room.
- Can't accommodate! $6 > 5$

- Infinite sets behave differently
- Suppose now the hotel has infinitely many rooms, indexed by N .

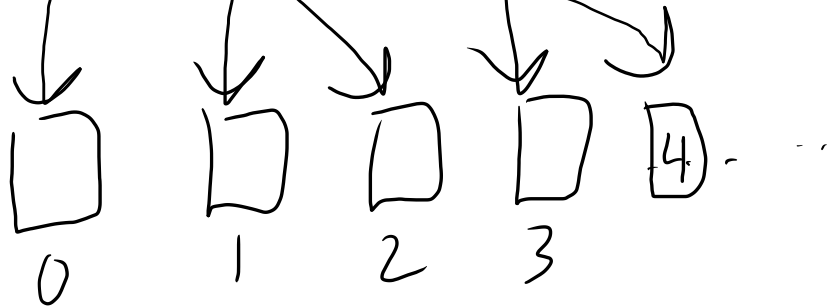


- Suppose all occupied, but additional person shows up demanding a room.
- Put him in room 0, everyone else moves one room to the right: $i \rightarrow i+1$.

Hotel $\mathbb{Z}$:



Hotel $\mathbb{N}$:



Can I evacuate $\mathbb{Z}$ to $\mathbb{N}$?

$\mathbb{N}$	$\mathbb{Z}$
0	0
1	1
2	-1
3	2
4	-2
5	3
6	-3
$\vdots$	

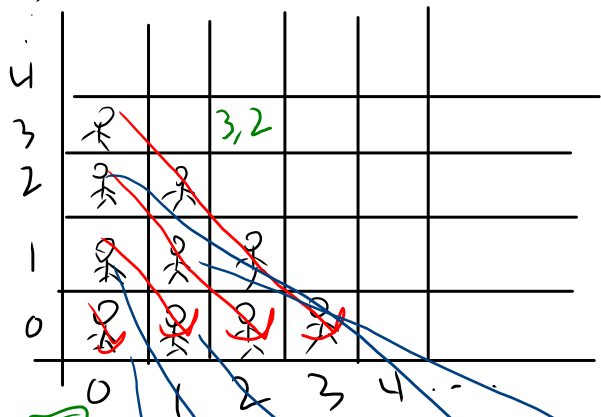
this is a one-to-one
correspondence between $\mathbb{Z}$ and $\mathbb{N}$.
(a.k.a. bijection)

So they are the same
'infinity'.

Hotel $\mathbb{N} \times \mathbb{N}$:

numerator
↓

denominator
→



fully occupied

$\frac{3}{2}$ represents $1\frac{1}{2}$

evacuate
↓

Hotel $\mathbb{N}$



...

Adding one, doubling, squaring doesn't change ∞ .

$$|\mathbb{N}| = |\mathbb{Z}| = |\mathbb{N} \times \mathbb{N}| = |\mathbb{Q}|$$

- All of these are the "same" infinity.
- we call these sets countably infinite

More generally:

- Two sets have the same cardinality iff there is a bijection between them

$$|A| = |B| \text{ iff } \exists \text{ bijection } f: A \rightarrow B$$

- $|A| = k$ ^{$k \in \mathbb{N}$} iff $|A| = |\{1, 2, \dots, k\}|$
iff $\exists \text{ bijection } f: A \rightarrow \{1, \dots, k\}$

- A is finite if $\exists k \in \mathbb{N}$ s.t. $|A| = k$

- A is infinite if it is not finite.

- we saw several infinite sets, but they all had same cardinality as $\mathbb{N}$.
- We call these countably infinite sets.
- Countable set: finite or countably infinite.

So far all infinite sets we saw were the same kind of infinity. I.e. they were all countably infinite.

Q: Is there a bigger infinity?
I.e. Is there a set which is not countable.

A: $\mathbb{R}$ is not countable.

$P(\mathbb{N})$ | / / / / ,
:
:

→ we will show this.

Thm: $\mathbb{R}$ are not countable.

Proof:

Represent $x \in \mathbb{R}$ in binary

$$2\frac{1}{4} \rightarrow 10.01_{\frac{1}{8}}$$

$$\frac{1}{16} \rightarrow 0.0001_{\frac{1}{16}}$$

$\uparrow \quad \uparrow \quad \uparrow$
 $\frac{1}{2} \quad \frac{1}{4} \quad \frac{1}{8}$

$$\frac{1}{3} \rightarrow 0.0101010101 \dots$$
$$\frac{1}{4} + \frac{1}{16} + \frac{1}{64} + \dots$$

- want to list all $\mathbb{R}$ in correspondence with $\mathbb{N}$.
- Make this even easier: Count real numbers > 0 and < 1

$0.1011011\dots$

- Assume for contradiction this is possible.

0	1	0	1	1	0	1				
1	0	1	0	1	0	1				
2	1	0	0	1	0	0	1	0	0	
3	1	0	1	1	1	0	1	1		
4	1	1	0	0	0	1	0	0	0	1
5										1
6										
7										

Diagonal

0.110101...

Diagonal Flipped

0.001010...

- The flipped diagonal disagrees with each row k on k th bit

- the k th flipped diagonal represents a real number nowhere in this list.

- Contradiction.



Book does same thing for $P(N)$

$$|P(N)| = |\mathbb{R}|$$

- We showed that $|\mathbb{R}|$ is a "bigger" infinity than $|\mathbb{N}|$.

Turns out: $|\text{Algorithms}| = |\mathbb{N}|$

$|\text{Problems}| = |\mathbb{R}| = |\mathcal{P}(\mathbb{N})|$

$|\text{Problems}| > |\text{Algorithms}|$

Halting problem is one specific example of a problem without algorithm to solve it.